

Note: Use of calculators, tablets, laptops, etc, is not allowed in the exam. Use of textbooks, lecture notes, or personal notes is not allowed either.

Note: If you have not completed both your computerized home assignments and your Coq mandatory assignment, you will not pass the course.

Total points: $6 + 16 + 15 + 20 + 12 + 1 = 70$.

1. (6 points)

Let A be the set $\{x, y, z\}$ and B be the set $\{x, y\}$.

- (a) Is A a subset of B ?
- (b) Is B a subset of A ?
- (c) What is $A \cup B$?
- (d) What is $A \cap B$?
- (e) What is $A \times B$?
- (f) What is the powerset of B ?

2. (16 points)

In all answers below, make sure that your automata are *complete*, i.e., that they have transitions defined from every state and for every input symbol.

- (a) Design a deterministic finite automaton that recognizes the language

$$\{w \in \{a, b\}^* \mid w \text{ starts with the substring } aab\}.$$

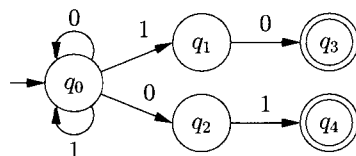
You can assume that the alphabet is $\{a, b\}$.

- (b) Design a deterministic finite automaton that recognizes the language

$$\{w \in \{a, b\}^* \mid w \text{ does not contain the substring } aab\}.$$

You can assume that the alphabet is $\{a, b\}$.

- (c) Let us consider the following non-deterministic automaton over the alphabet $\{0, 1\}$:



Describe the language recognized by the automaton verbally, in one or two sentences. Give a deterministic finite automaton that recognizes the same language.

3. (15 points)

(a) Give a regular expression that describes the language

$$L = \{w \in \{a, b\}^* \mid w \text{ begins and ends with the same symbol}\}$$

Note: the empty string does *not* belong to the above language.

(b) Consider the regular expression $(b \cup (aa^*b))^*$ over the alphabet $\{a, b\}$. Give the complete deterministic finite state automaton with *minimal number of states* that recognizes the language described by this regular expression.

(c) Give a regular expression that describes the language

$$L = \{w \in \{a, b\}^* \mid w \text{ does not contain the substring } aa\}$$

4. (20 points)

Let $G = (V, \Sigma, R, S)$ be a context-free grammar, where $V = \{S, T, U\}$, $\Sigma = \{0, \#\}$, and R is the set of rules:

$$S \rightarrow TT \mid U$$

$$T \rightarrow 0T \mid T0 \mid \#$$

$$U \rightarrow 0U00 \mid \#$$

- Can the string $\#$ be generated by the grammar? If so, give a parse tree for it. If not, provide explanation.
- Can the string $00\#$ be generated by the grammar? If so, give a parse tree for it. If not, provide explanation.
- Can the string $0\#\#0$ be generated by the grammar? If so, give a parse tree for it. If not, provide explanation.
- Give two more strings, different from all strings above and from each other, that can be generated by the grammar: show their derivations.
- Give two more strings, different from all strings above and from each other, that cannot be generated by the grammar. The strings must use the alphabet Σ above. Provide explanations why they cannot be generated by G .
- Prove that the language generated by G is not regular.
- Give a PDA (pushdown automaton) that recognizes the same language as that generated by G .

5. (12 points)

- Define the concepts “Turing-recognizable language” (also called “recursively enumerable language”) and “Turing-decidable language” (also called “recursive language”).

(b) Consider the language

$$\{x \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}^* \mid x \text{ is a product of two prime numbers}\}.$$

For instance, the string 15 is in the language as $3 \times 5 = 15$ but 16 is not in the language. Is the language Turing-decidable? Justify your answer.

(c) Prove the following claim either correct or incorrect: Let L_1 and L_2 be languages over an alphabet Σ . If L_1 is context-free and L_2 is Turing-decidable, then also the language $L_1 \cap L_2$ is context-free.

6. (1 point)

At what time did you finish answering the exam questions?

