

As in all examinations, you must answer the questions by yourself, without any outside help.

Using the material and software tools of the course is allowed.

Assignment 1 Formulas and normal forms. (Max. 5p)

Consider the propositional formula $a \vee (b \oplus c)$.

1. Give the truth table for the formula and all its subformulas.
2. Derive a logically equivalent formula that is in the conjunctive normal form.

Assignment 2 Conflict-Driven Clause Learning (Max. 5p)

Consider the CNF formula $(a \vee \neg f \vee \neg h) \wedge (b \vee c \vee \neg d) \wedge (c \vee \neg e) \wedge (d \vee e \vee f) \wedge (\neg f \vee \neg g) \wedge (g \vee h)$.

- (a) Simulate the conflict-driven clause learning (CDCL) satisfiability algorithm by choosing decision variables in the alphabetic order $a, b, c, \dots$, always setting the chosen variable to false, *until a conflict is reached*. Illustrate the resulting implication graph.
- (b) Using the constructed implication graph, illustrate the “first UIP cut” and give the corresponding learned clause.

Assignment 3 CSPs and Arc Consistency (Max. 5p) Consider the sum constraint $2x_1 + 2x_2 + x_3 + x_4 + 2x_5 = z$ when the domains are $D(x_1) = \{1, 2\}$, $D(x_2) = \{1, 3\}$, $D(x_3) = \{2, 3\}$, $D(x_4) = \{1, 3\}$, $D(x_5) = \{0, 1, 2\}$, and $D(z) = \{1, 3, 5, 9\}$.

Answer the following questions by using the dynamic programming approach described in the lecture slides. Illustrate the relevant parts of the graph presentation in your answer.

- Does the constraint have solutions? If it has, give one. If it does not, argue why this is the case.
- Is the constraint arc-consistent? If it is, explain why this is the case. If it is not, filter the constraint to be arc-consistent, and give the resulting new domains.

Assignment 4 Binary Decision Diagrams (Max. 5p)

1. Construct an OBDD for the Boolean function that yields *true* if and only if A and D have the same value and B and C have the same value. The variable ordering is A, B, C, D .
2. Annotate each node of your OBDD with the corresponding model-count.

Assignment 5 Modal logics (Max. 5p)

For each of the following formulas, give a model in which the formula is true in w_0 , or argue why the formula cannot be true in any state/world.

1. $a \wedge (\Box \neg a) \wedge (\Box \Box a)$
2. $(AGEF \neg a) \wedge EG a$

Assignment 6 (Max. 5p)

Explain why every possible directed graph, consisting of a set of $\leq 2^n$ states and a set of directed edges between states, can be expressed in terms of a set X of n Boolean state variables and a transition formula Φ over X and X' (where the latter has the “next state” variants x' of the state variables $x \in X$). How do you systematically construct Φ for any such given graph?

NB: Assignments continue on the other side of the question sheet!

Assignment 7 State-space search with propositional logic (Max. 5p)

Consider the following reachability problem in the propositional logic.

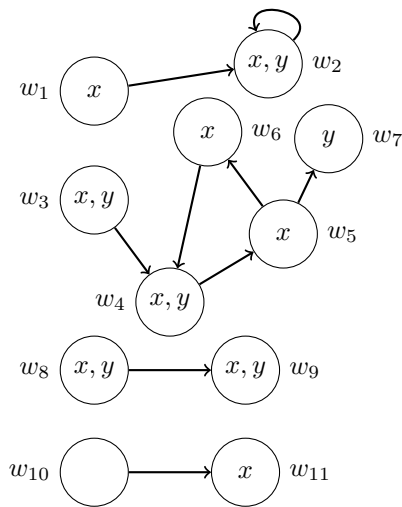
The states are encoded by the state variables a , b and c . The transitions are

- inverting the values of a and b and c (e.g. turning 101 to 010),
- inverting the values of a and c , and
- inverting the value of a .

Questions:

1. Give the formula for the transition relation that includes all 3 transitions.
2. Show how the image of the formula a w.r.t. the transition relation is computed.
3. How many steps does the symbolic breadth-first search algorithm run starting from $a \wedge \neg b \wedge c$? How many different states are reachable?

Assignment 8 (Max. 5p)



Run the CTL model-checking algorithm for the formula $\phi = EGEFy$ and the model on the left. For all of the relevant subformulas of ϕ , list which worlds will be labelled with that subformula.

Assignment 9 Satisfiability Modulo Theories (Max. 10p)

- (a) Recall that T_{EUF} is the theory of uninterpreted functions and predicates (with built-in equality, as usual) and T_{LIA} is the theory of linear arithmetic over integers. Which of the following are true? Justify your answer by giving a model or a proof.

1. $(f(x) \not\approx x) \rightarrow (f(f(x)) \not\approx x)$, where f is a unary function, is T_{EUF} -valid.
2. $(y \geq x) \wedge (y < x + 3) \wedge (y \geq -2x) \rightarrow (x \geq -1)$ is T_{LIA} -valid. (5p)

- (b) Show how the method presented in the material (constraint graphs and the Bellman–Ford algorithm) determines whether the following conjunction

$$(x - y \leq 3) \wedge (y - z \leq -1) \wedge (y - u \leq -2) \wedge (w - u \leq -3) \wedge (u - x \leq -1) \wedge (u - z \leq 2) \wedge (z - x \leq -2)$$

is *IDL*-satisfiable, where *IDL* is the integer difference logic. If the conjunction is *IDL*-unsatisfiable, find a negative cycle and give the corresponding *IDL*-unsatisfiable sub-conjunction. If it is *IDL*-satisfiable, give a model for the conjunction. (5p)

Assignment 10 At what time did you complete answering the questions?