

Please note the following: your answers will be graded only if you have passed all the three home assignments before the exam!

Assignment 1 (10p)

- (a) Define the following concepts: *modus ponens*, *disagreement set*, and *free variable occurrence*. (3 × 2p)
- (b) What is meant by the notation $\Sigma \models \phi$?
Prove in detail that if $\Sigma \models \phi$ and $\Sigma \models \neg\phi$ for some sentence ϕ , then Σ is unsatisfiable. (4p)

Assignment 2 (10p) Prove the following claims using semantic tableaux:

- (a) $\models (A \rightarrow B) \wedge (\neg A \rightarrow \neg B) \rightarrow (A \wedge \neg B \leftrightarrow B \wedge \neg A)$.
- (b) $\models \forall x(\exists yR(x,y) \rightarrow P(x)) \rightarrow \forall y\forall x(P(x) \vee \neg R(x,y))$.

Tableau proofs must contain all intermediary steps !!!

Assignment 3 (10p) Derive a Prenex normal form and a clausal form (i.e. a set of clauses S) for the sentence

$$\neg(\exists x(P(x) \vee \forall yQ(x,y)) \rightarrow \exists y(P(y) \vee Q(y,y))).$$

Make S as simple as possible. Prove that S is unsatisfiable using resolution.

Assignment 4 (10p) Let us represent strings “”, “a”, “b”, “aa”, “ab”, “ba”, “bb”, ... that consist of letters a ja b using ground terms

$$e, a(e), b(e), a(a(e)), a(b(e)), b(a(e)), b(b(e)), \dots,$$

built of a constant symbol e , which represents the empty string “”, and unary functions $a(x)$ and $b(x)$, that append the respective letter a or b at the beginning of a string x . Thus $a(b(e))$ is interpreted as $a(b(\text{""})) = a(\text{“b”}) = \text{“ab”}$.

- (a) Define predicate $U(x) = \text{“the string } x \text{ has either occurrences of the letter } a \text{ only or occurrences of the letter } b \text{ only but not both”}$ using predicate logic so that your definition covers all finite strings as described above.
- (b) Give a model $\mathcal{S} \models \Sigma$ of your definition Σ on the basis of which it holds that

$$\Sigma \not\models U(e).$$

Assignment 5 (10p)

Explain how the *weakest precondition* B_1 of an if-statement

$$\text{if}(B) \text{ then } \{C_1\} \text{ else } \{C_2\}$$

can be formed given a postcondition B_2 for it.

Consider the following program Double:

$$v=0 ; z=x ; \text{while}(! (z==0)) \{z=z-1 ; v=v+2\}.$$

Use weakest preconditions and a suitable invariant to establish

$$\models_p [\text{true}] \text{Double} [v==2 * x].$$

The name of the course, the course code, the date, your name, your student number, and your signature must appear on every sheet of your answers.